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Surface Area is roughly at least the volume!

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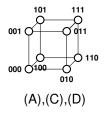
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101 〇 001 〇	111 O O 011	
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(A),(C),(D)		

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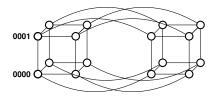
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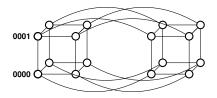
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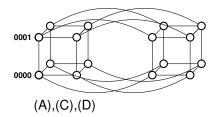
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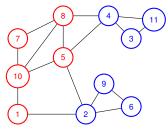
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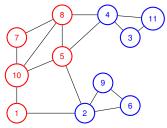
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Restatement: for any cut in the hypercube, the number of cut edges is at least the size of the small side.

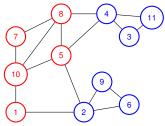


S is red, V - S is blue.



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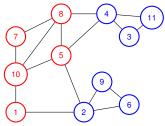
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What is size of cut?

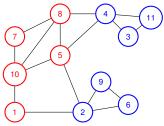
Number of edges between red and blue.



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Number of edges between red and blue. 4.



S is red, V - S is blue.

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Number of edges between red and blue. 4.

Hypercube: any cut that cuts off *x* nodes has $\ge x$ edges.

Proof of Large Cuts.

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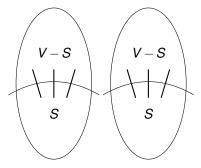
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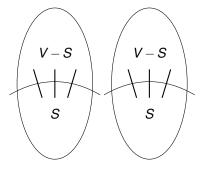
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Case 2: Count inside and across.

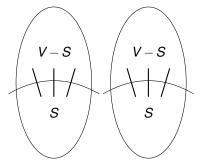


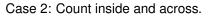
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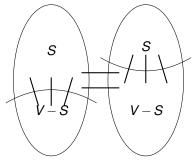
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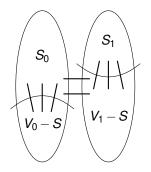
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$$|S_0| \ge |V_0|/2.$$

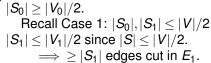


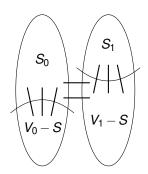
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 S_0 $V_0 - S$ $V_1 - S$ $|S_0| \ge |V_0|/2.$ Recall Case 1: $|S_0|, |S_1| \le |V|/2$ $|S_1| \le |V_1|/2$ since $|S| \le |V|/2.$

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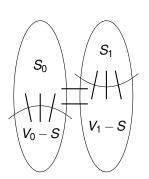
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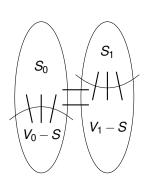
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$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| &\leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| &\leq |V|/2. \\ &\implies &\geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| &\leq |V_0|/2 \end{split}$$

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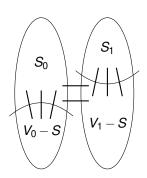
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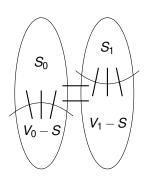


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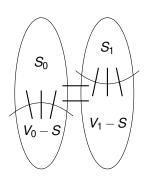


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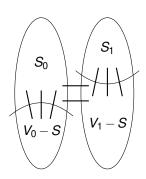


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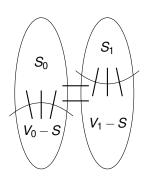
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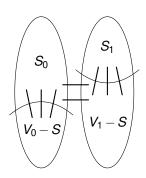
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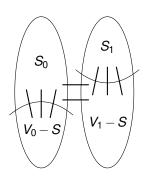
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Total edges cut:

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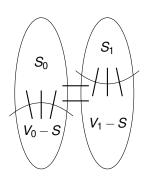
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 $\geq |S_1| + |V_0| - |S_0|$

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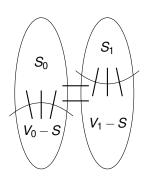
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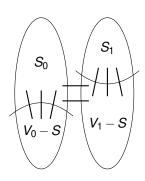
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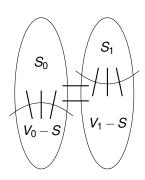
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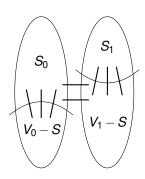
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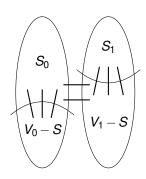
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 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0|$ $|V_0| = |V|/2 \geq |S|.$

Also, case 3 where $|S_1| \ge |V|/2$ is symmetric.

Hypercube proof: poll

Hypercube has large cuts proof uses these ideas:

- (A) If cuts are same size on two sides it works by induction.
- (B) Uses the fact that it is planar.
- (C) Recursive definition of hypercube.
- (D) If different size, can count edges between to subcubes.
- (E) Applies Euler's formula.

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(A),(D), and (E).

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Central object of study.

Modular Arithmetic.

Applications: cryptography, error correction.

Theorem: If d|x and d|y, then d|(y-x).

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Theorem: Every number $n \ge 2$ can be represented as a product of primes.

Proof: Either prime, or $n = a \times b$, and use strong induction. (Uniqueness? Later.)

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
- (B) Euler's formula.
- (C) The definition of a prime number.
- (D) Euclid's Lemma.

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
- (B) Euler's formula.
- (C) The definition of a prime number.
- (D) Euclid's Lemma.
- (A) and (C)

Next Up.

Modular Arithmetic.

If it is 1:00 now.

If it is 1:00 now. What time is it in 2 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

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What time is it in 100 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

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What time is it in 100 hours? 101:00!

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What time is it in 100 hours? 101:00! or 5:00.

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16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00. $101 = 12 \times 8 + 5$.

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What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

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Clock time equivalent up to addition of any integer multiple of 12.

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Custom is only to use the representative in $\{12, 1, \dots, 11\}$

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 $101 = 12 \times 8 + 5.$

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Custom is only to use the representative in $\{12, 1, ..., 11\}$ (Almost remainder, except for 12 and 0 are equivalent.)

This is Thursday is September 16, 2021.

This is Thursday is September 16, 2021. What day is it a year from then?

This is Thursday is September 16, 2021. What day is it a year from then? on September 16, 2022?

This is Thursday is September 16, 2021. What day is it a year from then? on September 16, 2022? Number days.

This is Thursday is September 16, 2021.

What day is it a year from then? on September 16, 2022? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

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Today: day 4.

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Today: day 4.

5 days from then.

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5 days from then. day 9

This is Thursday is September 16, 2021.What day is it a year from then? on September 16, 2022?Number days.0 for Sunday, 1 for Monday, ..., 6 for Saturday.

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This is Thursday is September 16, 2021. What day is it a year from then? on September 16, 2022? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

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5 days from then. day 9 or day 2 or Tuesday.

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What day is it a year from then? on September 16, 2022?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.

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5 days from then. day 9 or day 2 or Tuesday.

25 days from then.

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two days are equivalent up to addition/subtraction of multiple of 7.

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What day is it a year from then?

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What day is it a year from then? Next year is not a leap year.

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5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

What day is it a year from then?

Next year is not a leap year. So 365 days from then.

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5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369.

This is Thursday is September 16, 2021.

What day is it a year from then? on September 16, 2022? Number days.

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subtract 7 until smaller than 7.

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Smallest representation:

subtract 7 until smaller than 7. divide and get remainder.

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369/7

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subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5.

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or September 16, 2022 is a Friday.

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or September 16, 2022 is a Friday.

80 years?

80 years? 20 leap years.

80 years? 20 leap years. 366×20 days

80 years? 20 leap years. 366×20 days 60 regular years.

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Hmm.

What is remainder of 366 when dividing by 7?

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Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

```
Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7?

```
Years and years...
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Get Day: $4 + 2 \times 20 + 1 \times 60$

```
Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$

```
Years and years...
```

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7?

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
60 regular years. 365 \times 60 days
Today is day 4.
It is day 4 + 366 \times 20 + 365 \times 60. Equivalent to?
```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7
```

```
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Today is day 4.
Get Day: 4 + 2 \times 20 + 1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
```

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Years and years...
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Get Day: 4+2 \times 20+1 \times 60 = 104
Remainder when dividing by 7? 104 = 14 \times 7 + 6.
Or February 11, 2101 is Saturday!
```

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Remainder when dividing by 7? 104 = 14 \times 7 + 6.
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```

Further Simplify Calculation:

```
Years and years...
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Remainder when dividing by 7? 104 = 14 \times 7 + 6.
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Or February 11, 2101 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

```
Years and years...
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Remainder when dividing by 7? $104 = 14 \times 7 + 6$. Or February 11, 2101 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $2+2\times 6+1\times 4=18$.

```
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Further Simplify Calculation:
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20 has remainder 6 when divided by 7.
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```

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Further Simplify Calculation:

20 has remainder 6 when divided by 7.

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Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 6. September 16, 2101 is Saturday.

"Reduce" at any time in calculation!

x is congruent to y modulo m or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by m.

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... or x and y have the same remainder w.r.t. m.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Mod 7 equivalence classes:

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woo / equivalence classes

 $\{\ldots, -7, 0, 7, 14, \ldots\}$

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Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\}$

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Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

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 $\{\ldots, -7, 0, 7, 14, \ldots\}$ $\{\ldots, -6, 1, 8, 15, \ldots\}$...

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or " $a \equiv c \pmod{m}$ and $b \equiv d \pmod{m}$

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 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$

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$$a \equiv c \pmod{m}$$
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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore,

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

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Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m

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Can calculate with representative in $\{0, \ldots, m-1\}$.

x (mod m) or mod(x,m)

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

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 $mod(x,m) = x - \lfloor \frac{x}{m} \rfloor m$

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by $m in \{0, ..., m-1\}$.

mod $(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

 $mod (29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12$

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12

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6 ≡

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Check! $4(3) = 12 = 5 \pmod{7}$.

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"Common factor of 4"

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"Common factor of 4" \implies 8*k* - 12*l* is a multiple of four for any *l* and *k* \implies

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For 8 modulo 12: no multiplicative inverse!

"Common factor of 4" \implies 8k - 12l is a multiple of four for any l and k \implies 8k \neq 1 (mod 12) for any k.

Poll

Mark true statements.

(A) Mutliplicative inverse of 2 mod 5 is 3 mod 5.

- (B) The multiplicative inverse of $((n-1) \pmod{n} = ((n-1) \pmod{n})$.
- (C) Multiplicative inverse of 2 mod 5 is 0.5.
- (D) Multiplicative inverse of $4 = -1 \pmod{5}$.
- (E) (-1)x(-1) = 1. Woohoo.

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(F) Multiplicative inverse of 4 mod 5 is 4 mod 5.

(C) is false. 0.5 has no meaning in arithmetic modulo 5.

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If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

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 \implies Prime factorization of *m* and *x* do not contain common primes.

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⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

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Greatest Common Divisor and Inverses.

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If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

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For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$ Two solutions! $x = 2,5 \pmod{6}$

Very different for elements with inverses.

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Not a bijection.

Poll

Which is bijection?

(A) f(x) = x for domain and range being \mathbb{R} (B) $f(x) = ax \pmod{(n)}$ for $x \in \{0, ..., n-1\}$ and gcd(a, n) = 2(C) $f(x) = ax \pmod{n}$ for $x \in \{0, ..., n-1\}$ and gcd(a, n) = 1

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Next up.

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Next up. Euclid's Algorithm.

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Euclid's Algorithm. Runtime.

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Therefore $d \mod (x, y)$.

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GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)). **Proof:** *x* and *y* have **same** set of common divisors as *x* and mod(x, y) by Lemma 1 and 2. □ish.

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Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

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Therefore $d \mod (x, y)$. And d | y since it is in condition.

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