# Today.





Secret Sharing.



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Correcting for loss or even corruption.

Share secret among *n* people.

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Two points make a line. Lots of lines go through one point.

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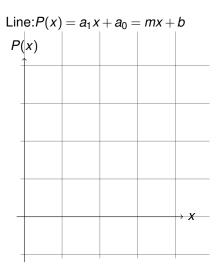
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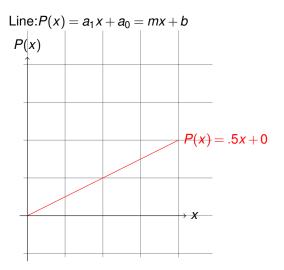
Polynomials P(x) with arithmetic modulo p: <sup>1</sup>  $a_i \in \{0, ..., p-1\}$  and

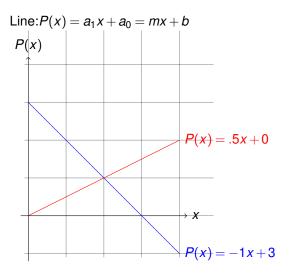
$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$
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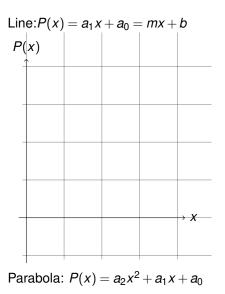
Line:  $P(x) = a_1 x + a_0$ 

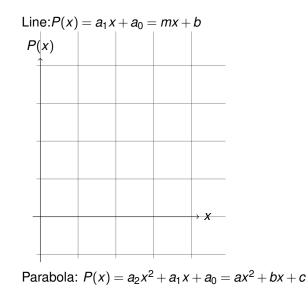
Line: $P(x) = a_1x + a_0 = mx + b$ 

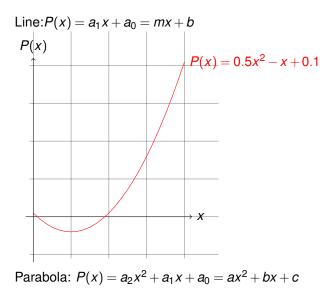


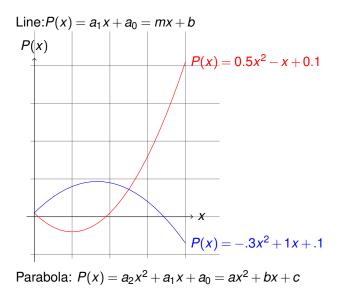


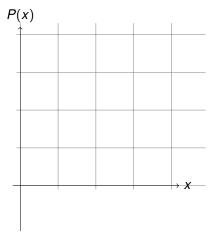


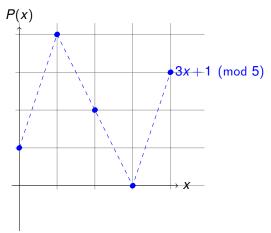


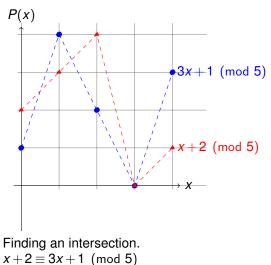




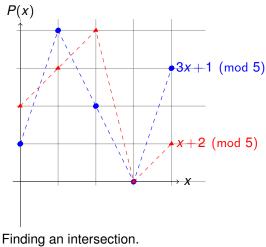




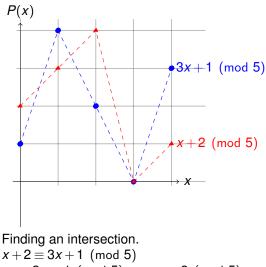




 $\implies 2x \equiv 1 \pmod{5}$ 



 $x + 2 \equiv 3x + 1 \pmod{5}$  $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5.



 $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Good when modulus is prime!! Two points make a line.

**Fact:** Exactly 1 degree  $\leq d$  polynomial contains d + 1 points.<sup>2</sup>

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**Modular Arithmetic Fact:** Exactly 1 degree  $\leq d$  polynomial with arithmetic modulo prime p contains d + 1 pts.

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### Poll.

#### Two points determine a line. What facts below tell you this?

Say points are  $(x_1, y_1), (x_2, y_2)$ .

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(A) Line is y = mx + b.

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All true.



Why solution? Why unique?

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#### Why solution? Why unique?

- (A) Solution cuz:  $m = (y_2 y_1)/(x_2 x_1), b = y_1 m(x_1)$
- (B) Unique cuz, only one line goes through two points.
- (C) Try:  $(m'x + b') (mx + b) = (m' m)x + (b b') = ax + c \neq 0$ .
- (D) Either  $ax_1 + c \neq 0$  or  $ax_2 + c \neq 0$ .
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Flow poll. (All true. (B) is not a proof, it is restatement.)

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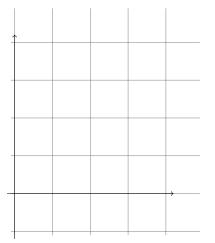
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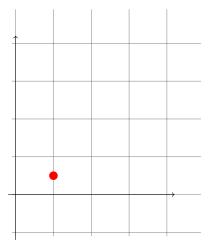
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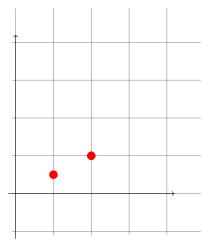
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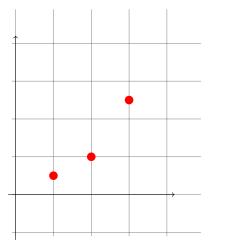
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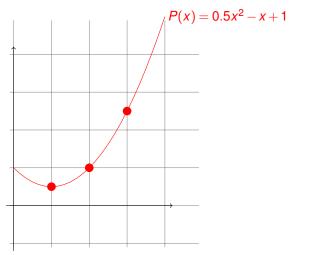
(A) and (D)

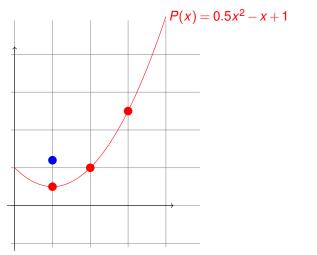


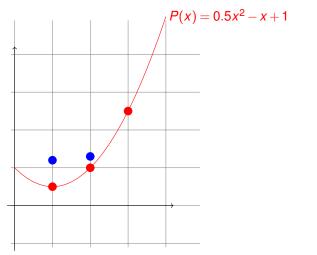


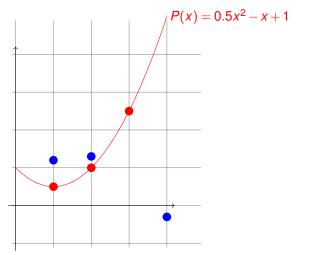


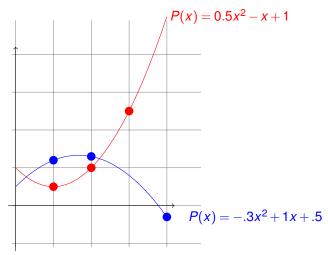




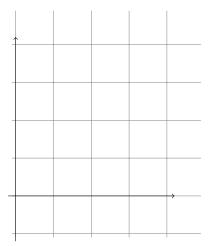


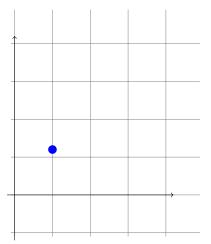


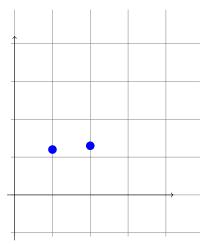


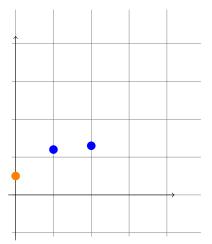


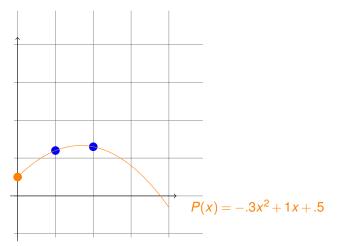
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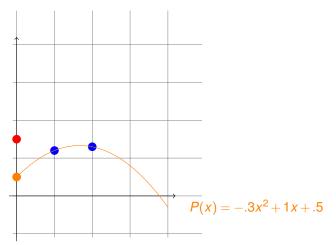


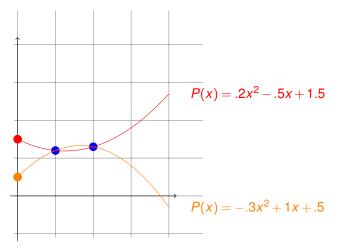


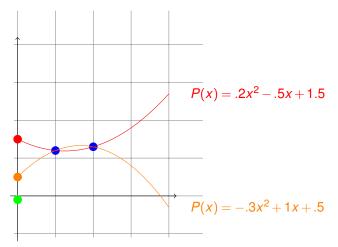


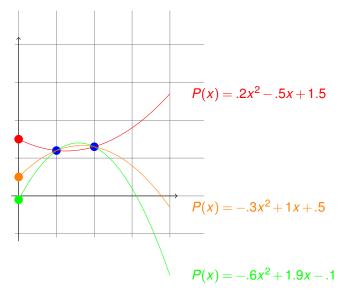


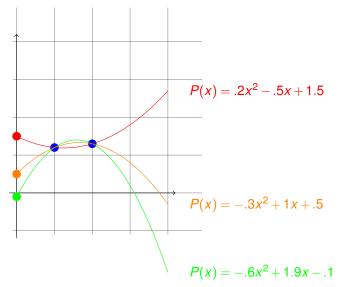












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# Poll:example.

The polynomial from the scheme:  $P(x) = 2x^2 + 1x + 3 \pmod{5}$ . What is true for the secret sharing scheme using P(x)?

(A) The secret is "2". (B) The secret is "3". (C) A share could be (1,5) cuz P(1) = 5(D) A share could be (2,4)(E) A share could be (0,3)

For a line,  $a_1x + a_0 = mx + b$  contains points (1,3) and (2,4).

*P*(1) =

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Backsolve:  $b \equiv 2 \pmod{5}$ . Secret is 2. And the line is...

 $x+2 \mod 5$ .

For a quadratic polynomial,  $a_2x^2 + a_1x + a_0$  hits (1,2); (2,4); (3,0).

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So polynomial is  $2x^2 + 1x + 4 \pmod{5}$ 

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**Modular Arithmetic Fact:** Exactly 1 degree  $\leq d$  polynomial with arithmetic modulo prime *p* contains *d* + 1 pts.

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Multiplicative inverses due to gcd(x,p) = 1, forall  $x \in \{1, ..., p-1\}$ 

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Mark what's true.

Poll

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(A) 
$$\Delta_1(x_1) = y_1$$
  
(B)  $\Delta_1(x_1) = 1$   
(C)  $\Delta_1(x_2) = 0$   
(D)  $\Delta_1(x_3) = 1$   
(E)  $\Delta_1(x_2) = 1$   
(F)  $\Delta_2(x_1) = 0$ 

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Find  $\Delta_1(x)$  polynomial contains (1, 1); (2, 0); (3, 0).  $\Delta_1(x) = \frac{(x-2)(x-3)}{(1-2)(1-3)} = \frac{(x-2)(x-3)}{2} = (2)^{-1}(x-2)(x-3) = 3(x-2)(x-3)$  $= 3x^2 + 3 \pmod{5}$ 

Put the delta functions together.

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Construction proves the existence of the polynomial!

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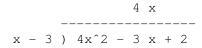
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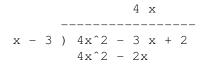
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In general, divide  $P(x)$  by  $(x - a)$  gives  $Q(x)$  and remainder  $r$ .  
That is,  $P(x) = (x - a)Q(x) + r$ 

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(Almost) the same as what is missing: one P(i).

## Runtime.

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Runtime: polynomial in k, n, and  $\log p$ .

- 1. Evaluate degree k 1 polynomial *n* times using log *p*-bit numbers.
- 2. Reconstruct secret by solving system of *k* equations using log *p*-bit arithmetic.

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Infinite number for reals, rationals, complex numbers!

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Compute solution: *m*,*b*.

Unique:

Assume two solutions, show they are the same.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root. Induction, says only *d* roots.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root.

Induction, says only *d* roots.

Apply: P(x), Q(x) degree d.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root.

Induction, says only *d* roots.

Apply: P(x), Q(x) degree d.

P(x) - Q(x) is degree  $d \implies d$  roots.

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root.

Induction, says only *d* roots.

Apply: P(x), Q(x) degree d.

P(x) - Q(x) is degree  $d \implies d$  roots.

P(x) = Q(x) on d+1 points  $\implies P(x) = Q(x)$ .

Two points make a line.

Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Cuz:

Solution: lagrange interpolation.

Unique:

Roots fact: Factoring sez (x - r) is root.

Induction, says only *d* roots.

Apply: P(x), Q(x) degree d.

P(x) - Q(x) is degree  $d \implies d$  roots.

P(x) = Q(x) on d+1 points  $\implies P(x) = Q(x)$ .

Secret Sharing:

k points on degree k - 1 polynomial is great!

Can hand out *n* points on polynomial as shares.